Collecting primes with $p^2 - 1\,827$ -smooth OR Reduced sets for likely solutions to the \$620 problem

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In memory of my friend Red Alford

Probable Primes and Pseudoprimes

Fermat's Little Theorem:

- If p is prime, then $a^p \equiv a \mod p$, for all integers a.
- The converse doesn't always follow...
- $2^{341} \equiv 2 \mod 341$.
- A probable prime to the base a is a number n such that $a^n \equiv a \mod n$.
- A pseudoprime to the base a is a composite probable prime to the base a.
- (Sometimes called a Fermat pseudoprime.)
- There aren't that many pseudoprimes (compared to primes).

Fibonacci pseudoprimes

- A Fibonacci pseudoprime is a composite n such that $n|F_{n-(\frac{n}{5})}$, where F_k is the k^{th} Fibonacci number.
 - (Generalizes to Lucas sequences and Lucas pseudoprimes.)
- Pomerance, Selfridge and Wagstaff offer \$620 for a base-2 pseudoprime that is also a Lucas pseudoprime and is 2 or 3 mod 5. (Or a proof that none exists.)

Carl's Heuristic

- In 1984, Carl Pomerance gave a heuristic (a modification of Erdos' heuristic for Carmichael numbers) that said that there should be infinitely many $(\gg x^{1-\epsilon})$ solutions to the \$620 problem.
- The paper is available at
 - http://www.pseudoprime.com/pseudo.html
- We choose a set Q of primes less than a bound B. Let Q_1 be the subset of Q consisting of primes congruent to $1 \mod 4$ (excepting 5). Let Q_3 be the subset of Q consisting of primes congruent to $-1 \mod 4$. Then we search for primes $p \equiv 3,27 \mod 40$ with (p-1)/2 squarefree and consisting only of primes in Q_1 and (p+1)/4 squarefree and consisting only of primes in Q_3 .
 - Call this set of primes P.

Carl's Heuristic, cont.

- Let $M_1 = \prod_{q \in Q_1} q$, and $M_3 = \prod_{q \in Q_3} q$.
- Let P' be a subset of P, and let $n = \prod_{p \in P'} p$. Assume that n has an odd number of prime factors, and further that $n \equiv 1 \mod M_1$ and $n \equiv -1 \mod 4M_3$. Then $n \equiv 2$ or $3 \mod 5$, n is a (strong) Fermat pseudoprime to the base 2 and m is a Fibonacci pseudoprime. (In fact, n is also a Carmichael number.)
- Why is *n* a Fermat pseudoprime? For each *p*|*n p* ∈ *P*, so we have $p 1|2M_1$. Further, $2M_1|n 1$, by the assumptions on *n*. Therefore, p 1|n 1. Therefore $2^{n-1} \equiv 2^{(p-1)\frac{n-1}{p-1}} \equiv 1$.

n is a Fibonacci pseudoprime by a similar argument.

Why does *n* exist?

- We assume that all possible *n*s are randomly distributed $\mod 4M_1M_3$. (This is not accurate, but it is probably pessimistic.) If $2^{|P|} > \varphi(4M_1M_3)$, then there is likely an *n* in the appropriate congruence class. If $2^{|P|} > 2\varphi(4M_1M_3)$, there is likely such an *n* with an odd number of prime factors.
- We call such a set a "likely solution".
- In the mid-1990s, Red Alford and I presented at SERMON a set *P* with cardinality 2030, where $2\varphi(4M_1M_3) \approx 2^{1812}$.

How to find *n***?**

- Heuristics say 2^{218} solutions! How to find 1?
- Trim P to minimum possible set. (|P| = 1812.)
- Naive way: form all possible subproducts; check if they win. Work: 2^{|P|}.
- Less naive way: form all possible subproducts of odd cardinality. Work: $2^{|P|-1}$.
- Better way: categorize *P* into equally sized subsets *P*₁ and *P*₂. Compute all possible subproducts of *P*₁ mod $4M_1M_3$. Compute all possible subproducts of *P*₂. Sort the two lists together in a clever way; if you get any matches, you win! Work: $2^{|P|/2}$.
- Practical way...Work: 2⁴⁰. (Unfortunately not known to exist.)

Relax!

- Carl's conditions were very strict. You can relax a number of them and still get a solution to the \$620 problem (though perhaps not a Carmichael number). For example, you need $ord_2(p)|2M_1$, not necessarily $p-1|2M_1$ (though the latter implies the former).
- Similarly, you can look at the "Fibonacci order" instead of p + 1.
- Also, if you look at primes that are 3 mod 4 instead of 3 mod 8, you lose the *strong* pseudoprime, but you get more primes to choose from. (You have to add powers of 2 to M₁.)
- You don't need $(p^2 1)/8$ squarefree.
- You don't need to categorize primes by their value mod4; you can be smarter.

Chen/Greene

- In a 2003 paper, Chen and Greene develop each of these ideas.
- They find a likely set with 1241 elements.
- They use 70 primes in each of M₁ and M₃. (Red and I used 100 each.)
- They carefully assign primes to M_1 and M_3 to balance them out.
- This paper renewed my interest in reducing the size of likely sets.

Don't be smart, be dumb!

- The general method for finding primes is constructing p-1 to be smooth, then testing p for primality and $ord_f(p)$ for smoothness.
- (Alternatively, construct p + 1...)
- Cycle through k-subsets of Q_1 and/or Q_3 .
- It's almost as cheap to test p + 1 for M_1M_3 -smoothness as M_3 -smoothness.
- Not horrifically more expensive to cycle through k-subsets of Q than of Q_1 .
- Let Q_1 and Q_3 choose themselves.

Method to the madness

- Construct k-subsets of Q with B = 811 for small k.
- Randomly separate Q into Q_1 and Q_3 . Repeat. Keep highest cardinality set.
- Try switching primes back and forth between Q_1 and Q_3 ...up to 9 primes at a time. See if it improves.
- Construct k-subsets of Q_1 and Q_3 for slightly larger k.
- Try switching primes again.
- Don't have enough primes. Bump B to 827; search for primes p where p + 1 is a multiple of 821, 823, or 827.
- Find likely set of size 1182. (71 primes in Q_1 ; 72 in Q_2 .)
- Celebrate! Write talk.
- Re-do more systematically. Write paper as if you knew the correct bound to begin with. (To do.)

On the horizon

- Generate more primes?
- (Increase size of k in both steps.)
- (Include primes > 827 as long as they don't divide the order.)
- Search smarter?
- One idea: let the primes be nodes on a graph. Connect two primes if they are "compatible".
 - $(gcd(ord_2(q_1), ord_f(q_2))|2, etc.)$
 - Find maximal complete subgraph.
 - This is probably NP-complete...
- Other types (Perrin Q-pseudoprimes)